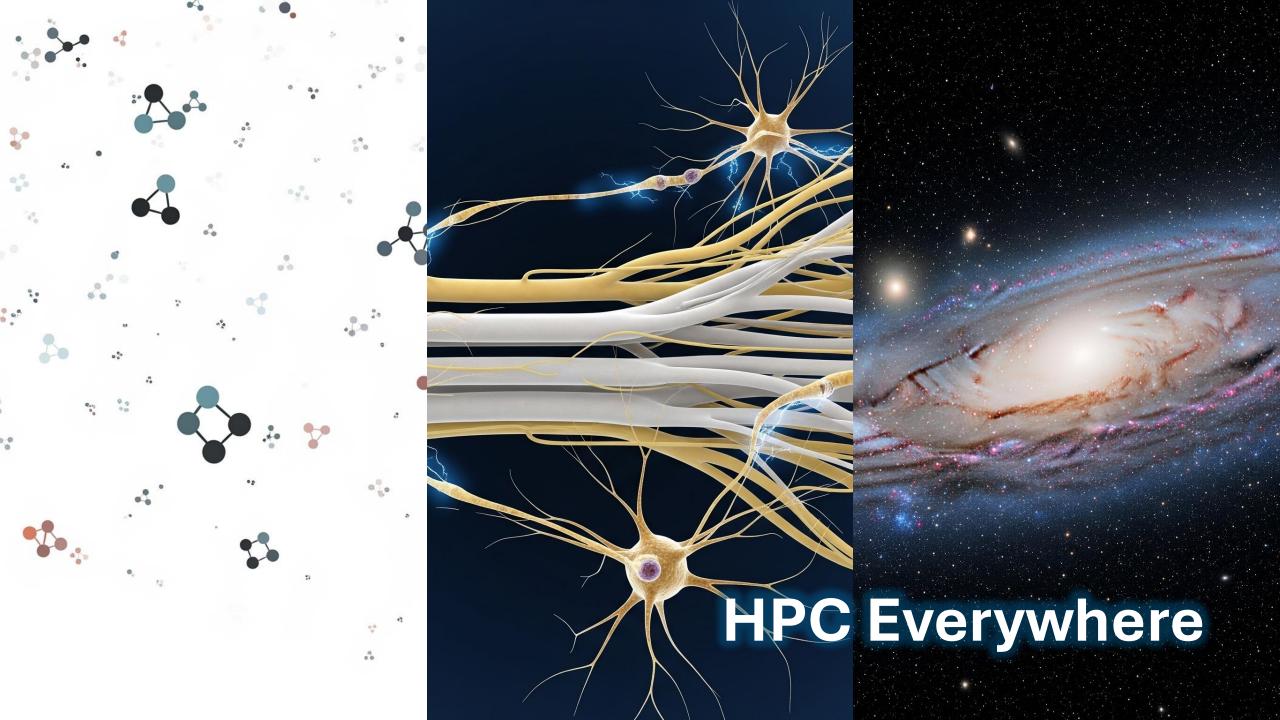
HPC Everyday & Everywhere

Kevin A. Brown, PhD

Assistant Computer Scientist Argonne National Laboratory kabrown@anl.gov





Kevin A. Brown

Researcher?

Research is my medium, not my passion.

Where my passion lies:

- Event Planning
- Program Development
- Community Building



Department of Entertainment (DOE)
Fun Opportunity Announcement (FOA)

Halloween Costume Party!

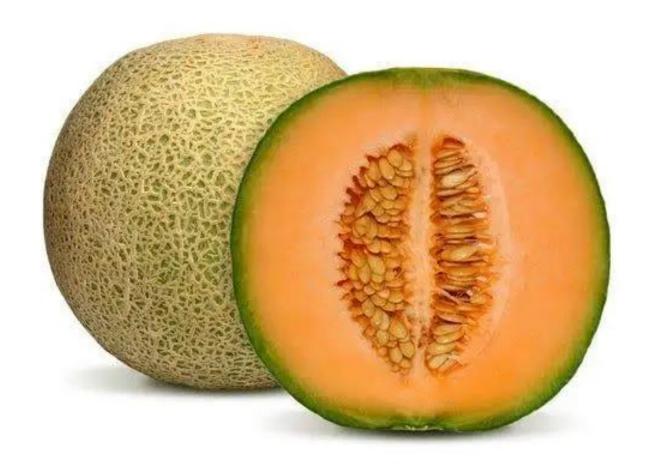
FOA Number: HW2024.10.25-B240

Application Deadline (party date):
October 25, 2024 at 4:30pm - 6:00pm

Follow-on fun (after party) is available starting 6:30pm @ ReNEW Westmont

Point of Contact (POC): Kevin B.



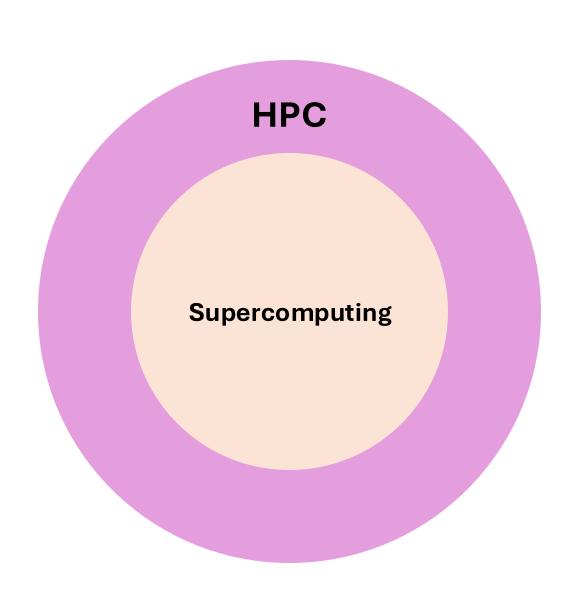


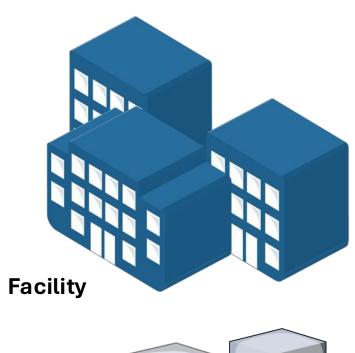


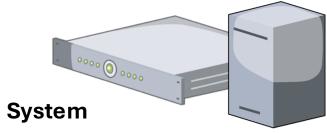
My Background

No Job Offer – PNNL, USA | English Teacher, Japan | Others MANY Rejected Papers – Various conferences **Not Selected -** Prime Minister Youth Award, Jamaica No Scholarship for 1st Year Undergrad - Jamaica Lost Scholarship in Final Year Undergrad - Jamaica Lost Friendships & Relationships – Jamaica, USA, Japan **Lost Election** – Student Union President, University of Technology, Jamaica

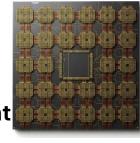
Major Research Mistakes – Too many to count







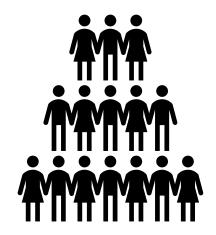




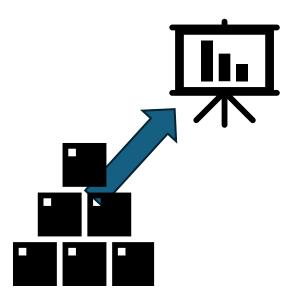
Infrastructure

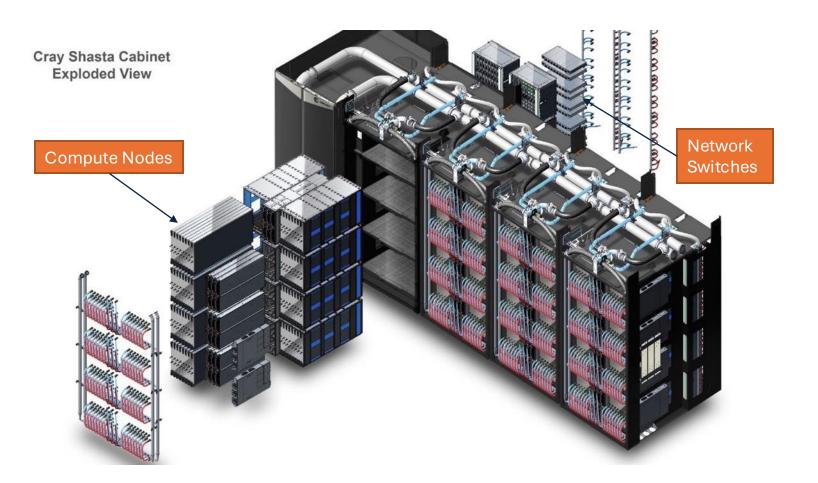


Application

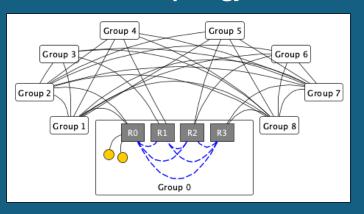


Goal

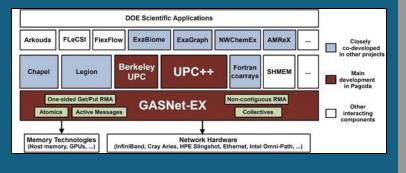




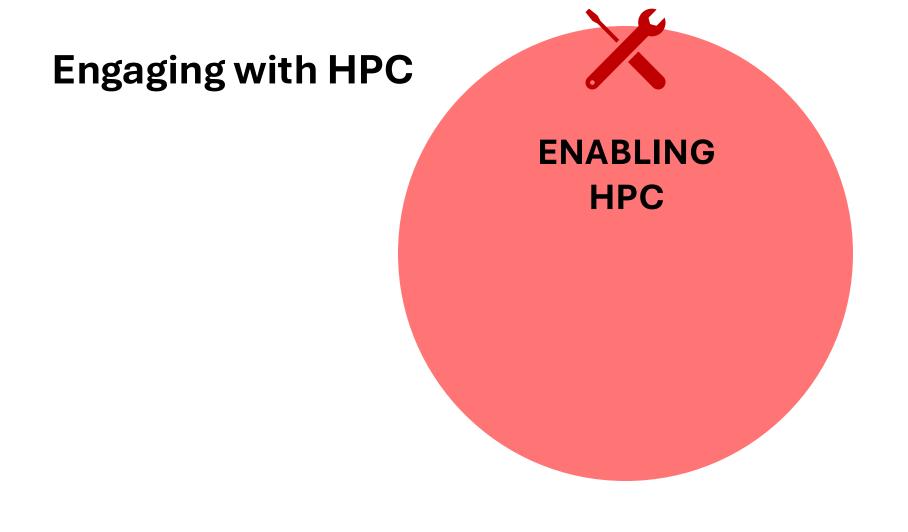
Network Topology View

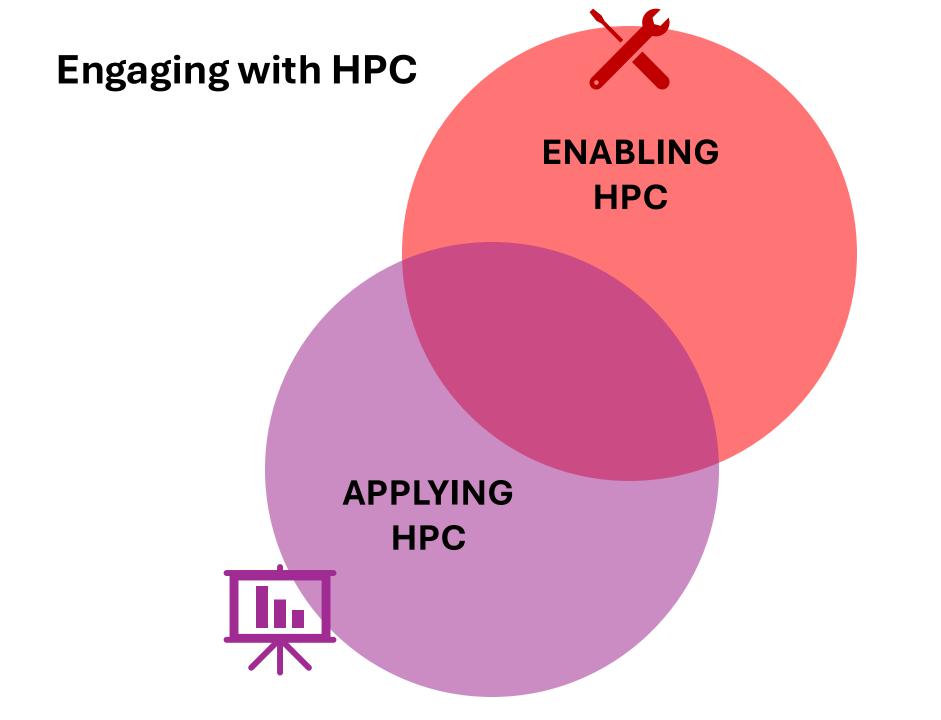


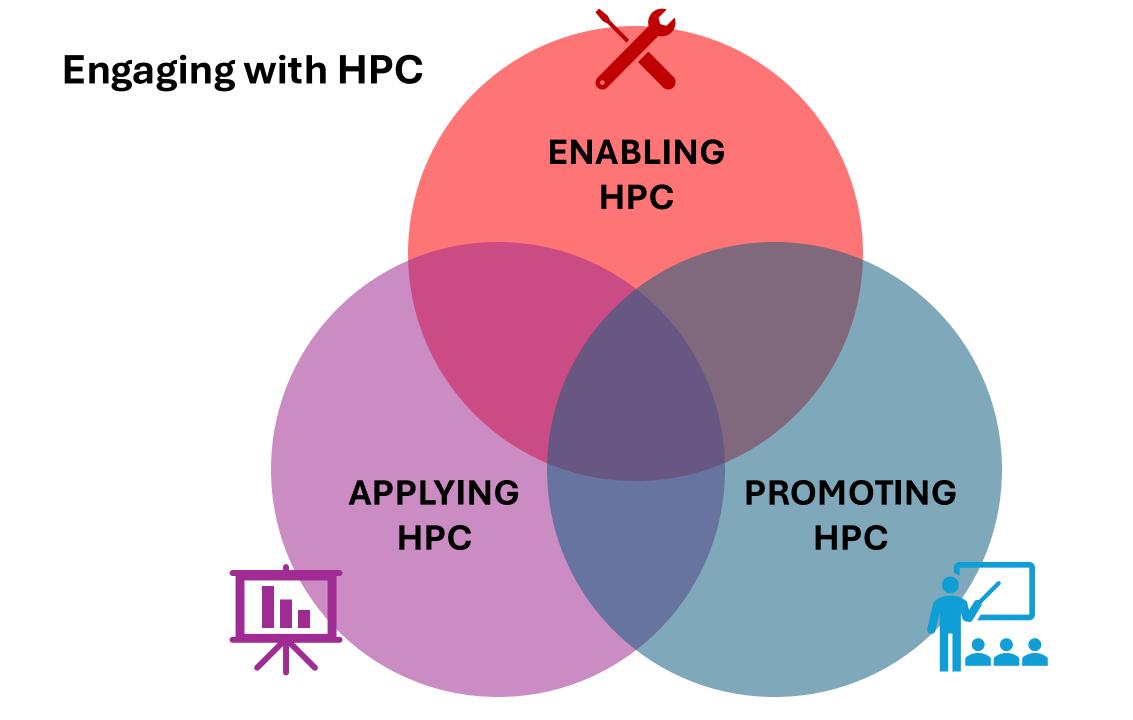
Software Stack View



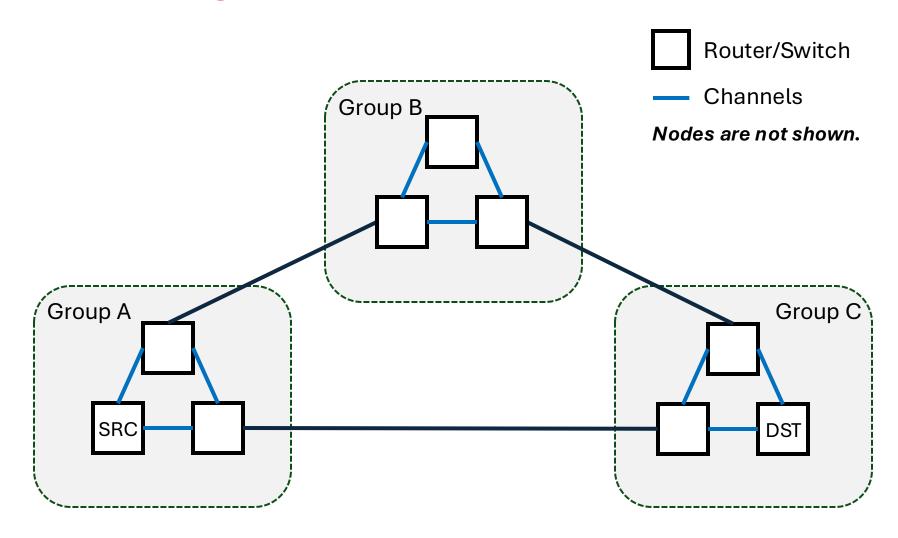




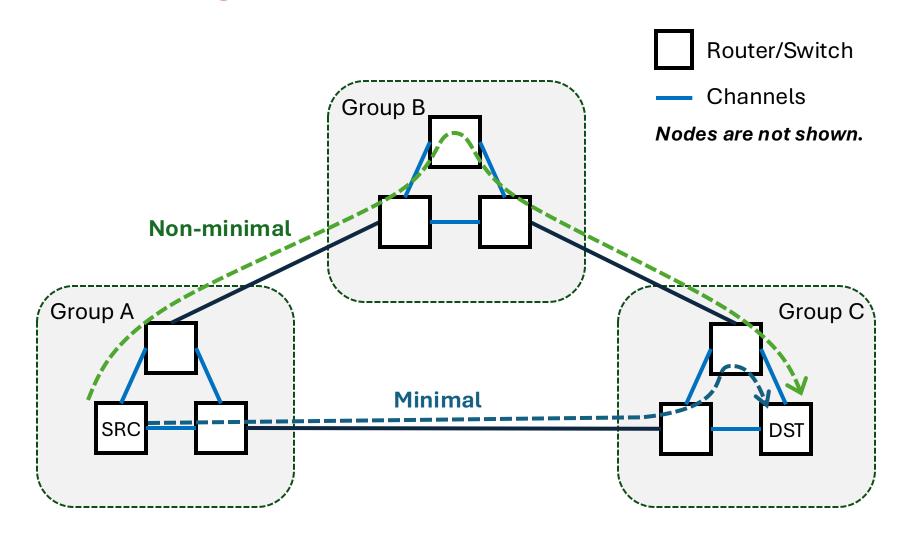




Network Topology

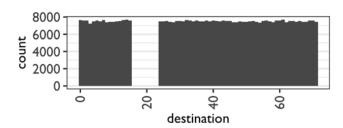


Network Topology



Performance Analysis

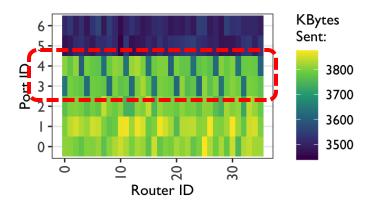
Count Messages Sent



Balanced distribution of packets



Total Traffic Per Switch Port



Unexpected traffic pattern across global ports



Understanding HPC Network Behavior Using Low-level Metrics

Kevin A. Brown and Robert B. Ross Mathematics and Computer Science Division, Argonne National Laboratory

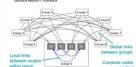
Supercomputer systems are complex interconnections of nodes supported by advanced networking. However, creating accurate simulation models is an error-prone process that requires accurately modeling real-world capabilities without full, complex technical re-implementations

To better evaluate network model activities, we use a performance refinement process to drill down into routing anomalies present in the dragonfly network model of the CODES simulation toolkit [1]. This refinement process starts with viewing high-level performance data and then iteratively uses lower-level, componentspecific data to locate the source of the anomaly.

We identify new low-level routing metrics that can expose the behavior of network model and attach these metrics to packets flowing through the system. Using these per-packet, per-hop routing metrics, we effectively attributed anomalies in system-wide traffic distribution to the routing algorithm's interaction with network connectivity configuration

DRAGONELY NETWORKS

- The 1D Dragonfly [2] is hierarchical network Each router is connected to a set of compute
- Routers are interconnected within logical
- Groups are then interconnected via "global"
- . There are multiple paths between source and



ADAPTIVE ROUTING

- Progressive adaptive routing (PAR) [3] attempts to use the fastest path between source and destination routers
- The congestion at potential paths is evaluated at each stop along the packet's journey
- The packet is re-routed along alternative nonminimal paths when the original paths is more



72-nodes, 9-group, 36-routers 1D Dragonfly network 2 nodes/router | 4 routers/group | 25Gb/s link bandwir

Uniform Global Random Traffic

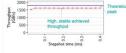
Simulated by the CODES/ROSS simulation toolkit [1]

· We use a uniform global random traffic pattern to evaluate PAR routing in CODES Traffic from each node is sent uniformly to all global destinations

Count Packets Sent by Nodes in Group 2



· Performance validation usually focuses on high-level behaviors, such as network



routing is achieving the expected performance FINDING:

· Port- and packet-specific metrics expose that a

source router with a minimal path

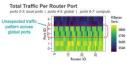
non-minimal path is rarely taken when it shares the

- The basic PAR algorithm in CODES compares

only the loads on min, and non-min, paths

- The PAR algorithm decide which port is used by a packet leaving each router
- Ports may become blocked due to congestion and cannot send data
- Online port-specific metrics show how well traffic load is balancing across ports

Total Traffic Per Router Port



Global Port Traffic For Router in Group 2



the traffic pattern caused by routing.

- . The intermediate group taken by a packet is recorded in each packet along its journey

 - This marks the non-min. path used
- . The load on non-minimal paths is confirmed to be unbalanced by viewing the path distributions

[Source Group = 2]

Illustration of Network Connectivity Highlighting Select Non-minimal and Minimal Paths



CONCLUSION

- . Taken alone, high-level metrics may not show anomalies in routing behavior
- Iteratively refining the analysis approach can systematically expose anomalies Component- and packet-specific measurements should be incorporated in the analysis
- The 2x minimal path bias will cause the non-min. Low-metric metrics requires storing large volumes path to appear more loaded of data -1ms of activity can generate GBs of data

NEXT STEPS

- Evaluate larger scale network behaviors Incorporate in-situ analytics framework for online
- streaming of data from the simulation Optimize the adaptive routing algorithm in CODES

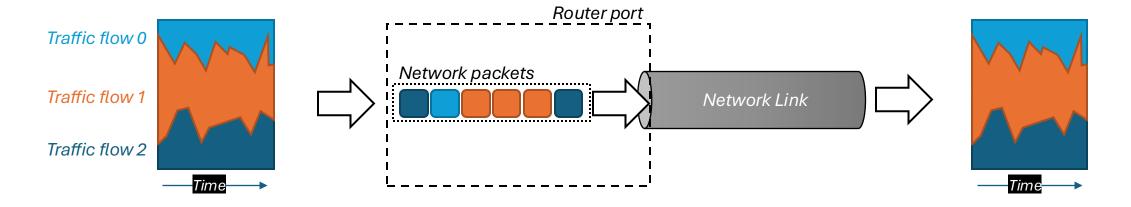


ENERGY U.S. Department of Energy laboratory managed by UChicago Argonne, LLC

Argonne 📤

Quality of Service

Network without QoS



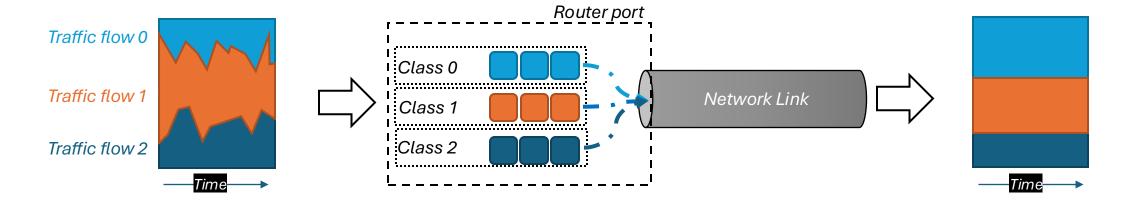
Variation in traffic load over time

No isolation; unregulated access to link

No performance guarantees

Quality of Service

Network with QoS

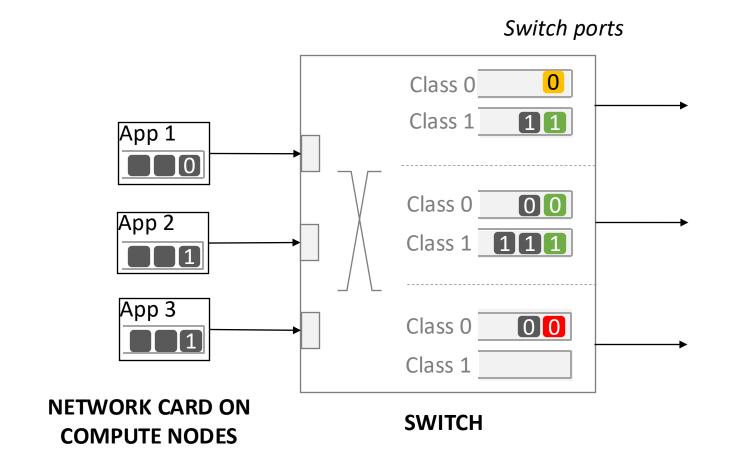


Variation in traffic load over time

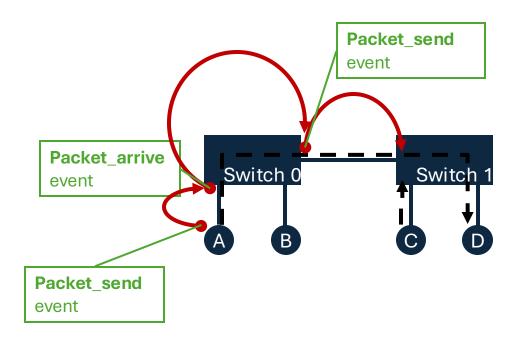
Isolation across multiple traffic classes; **QoS** policy regulates access to link

Consistent performance guaranteed

Quality of Service



Simulating Supercomputer Networks

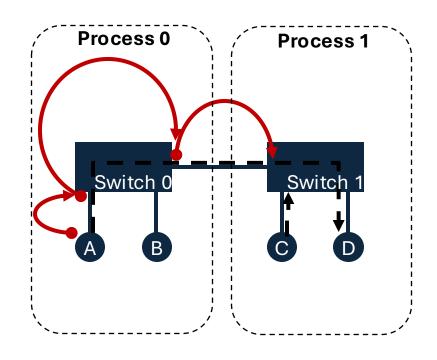


Network model

Time	Switch 0 events	Switch 1 events
5	Packet(A)Arrive*	-
10	Packet(A)Send	-
15	-	Packet(A)Arrive
16	-	Packet(C)Arrive*
20	-	Packet(A)Send
21	-	Packet(C)Send

Events in the network model

Simulating Supercomputer Networks



Network model

Process 0

Time	Switch 0 events
5	Packet(A)Arrive*
10	Packet(A)Send

Process 1

Time	Switch 1 events	
15	Packet(A)Arrive	
16	Packet(C)Arrive*	
20	Packet(A)Send	
21	Packet(C)Send	

Events in the network model

Teaching and Supporting HPC





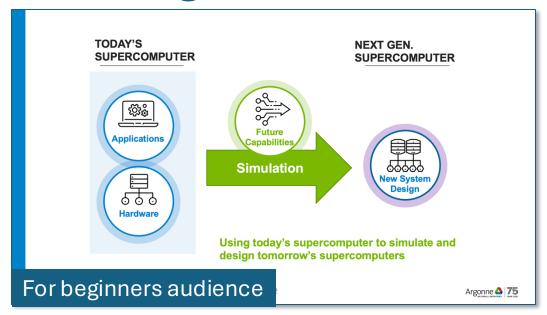
SC24 Early Career Program Chair

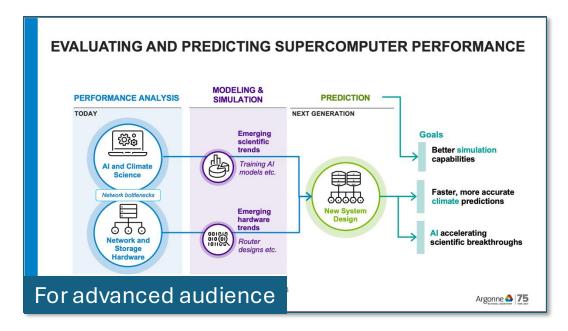
- Proposing Work
- Reporting Work
- Managing Work
- Career Paths Panel

CARLA 2025General Co-chair

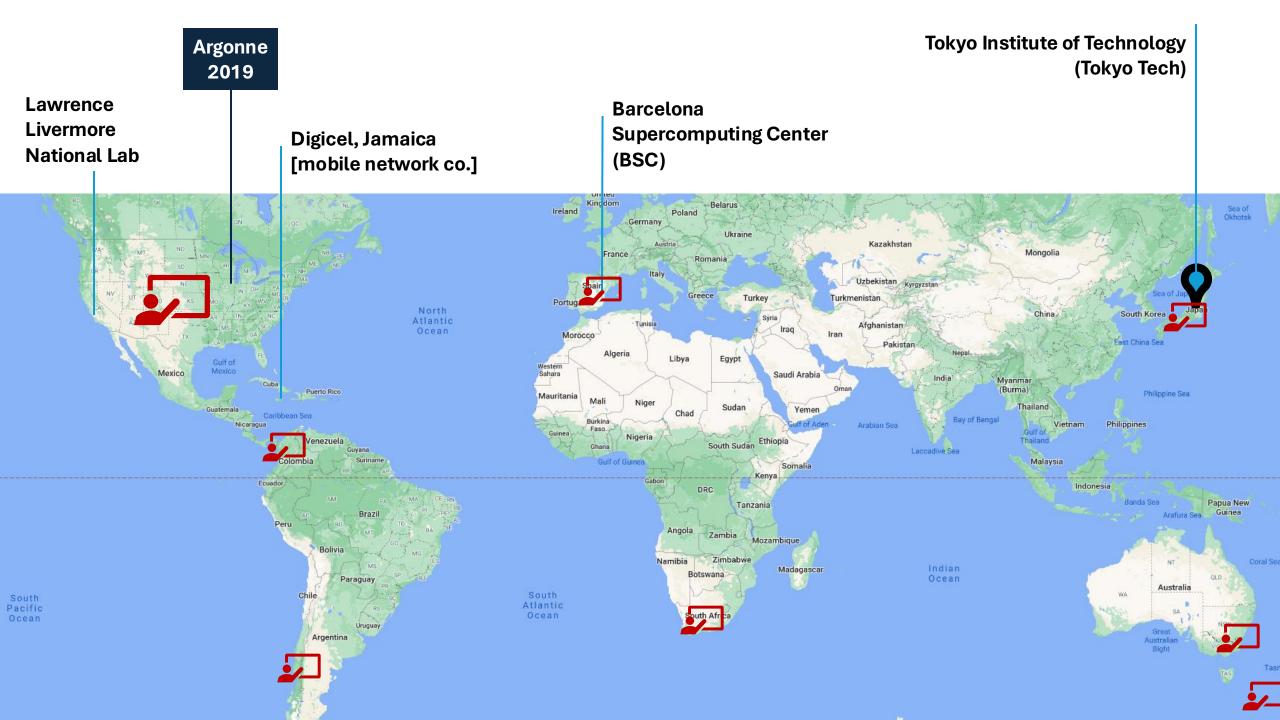
- Keynotes (NASA, Argonne Lab, etc.)
- HPC/AI workshops, tutorials, et al.
- Exhibitor booth
- Networking events

Promoting HPC

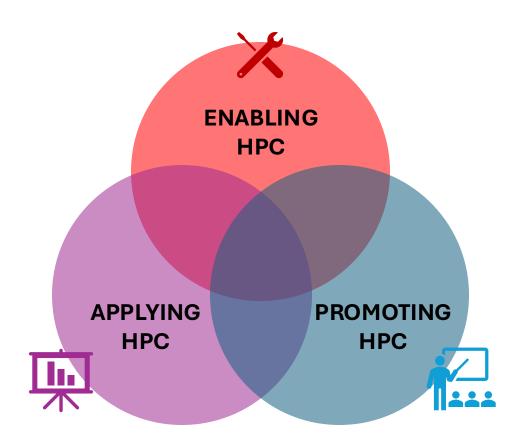








Engaging with HPC



The Journey

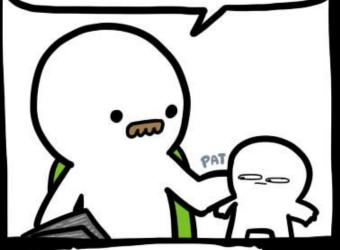
- 1. It's OK to not know what you want to do
 - It may not exist yet
- 2. Your path to HPC may not be linear
- 3. Your work with HPC may not be well bounded













SRGRAFO